# PIMS-MITACS Summer School on Quantum Information Science

Non locality

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#### Speed of communication

- « Faster than light communication is absolutely impossible»
- « The independent realisation of an experiment in a far away location should not instantaneously influence the outcome of an experiment in my lab»
- «The outcome of a measurement carried out on a system only depend on that system»

#### Pseudo-telepathy

« The capacity to perform a task that requires communication without signaling »

### Setup stage



## The game



#### The game

	Alice	Bob
Questions (1 bit)	$x_A \square \{A, B\}$	$x_B \square \{B,C\}$
Answer (1 bit)	$y_A \square \{0,1\}$	$y_B \square \{0,1\}$

Rule of the game: If  $x_A = x_B = B$  then  $y_A = y_B$ The game is to be repeated numerous times.

$$P_{XY} = \Pr[y_A = y_B \mid x_A = X \text{ and } x_B = Y]$$

The goal is to maximize  $P_{AB} + P_{BC} \square P_{AC}$ 

#### The strategy

Assume Alice and Bob have agreed on *any* strategy that does not involve communication. Randomness is allowed.

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In round i if x_A = A Alice outputs a_i.
In round i if x_A = B Alice outputs b_i.
In round i if x_B = B Bob outputs b_i.
In round i if x_B = C Bob outputs c_i.
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#### The inequality

$$\Pr[a = b = c] + \Pr[a = b \neq c] + \Pr[a \neq b = c] + \Pr[a \neq b \neq c] = 1$$

$$Pr[a = b] + Pr[b = c] \square Pr[a = c]$$

$$= (Pr[a = b = c] + Pr[a = b \neq c]) + (Pr[a = b = c] + Pr[a \neq b = c]) \square (Pr[a = b = c] + Pr[a \neq b \neq c])$$

$$= Pr[a = b = c] + Pr[a = b \neq c] + Pr[a \neq b = c] \square Pr[a \neq b \neq c]$$

$$\square Pr[a = b = c] + Pr[a = b \neq c] + Pr[a \neq b = c] + Pr[a \neq b \neq c]$$

$$= 1$$

$$Pr[a = b] + Pr[b = c] \square Pr[a = c] \square 1$$

$$P_{AB} + P_{BC} \square P_{AC} \square 1$$

#### The contradiction

Alice and Bob are separated (several lightyears away) and numerous instances of the test are performed. The collected data shows the following correlations:

$$P_{AB} = P_{BC} = \frac{3}{4}$$
  $P_{AC} = \frac{1}{4}$   $P_{AB} + P_{BC} \square P_{AC} = \frac{5}{4} > 1$ 

Pseudo-telepathy is then possible! Is something wrong?

For each round Alice and Bob share a precious pair of entangled qubits.

$$\left| \Box \right\rangle = \frac{1}{\sqrt{2}} \left( 00 \right)_{AB} + \left| 11 \right\rangle_{AB} \right)$$

$$U | 0 \rangle = \frac{\sqrt{3}}{2} | 0 \rangle + \frac{1}{2} | 1 \rangle$$

$$U | 1 \rangle = \frac{1}{2} | 0 \rangle \Box \frac{\sqrt{3}}{2} | 1 \rangle$$

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$$U | 1 \rangle = \frac{1}{2} | 0 \rangle \Box \frac{\sqrt{3}}{2} | 1 \rangle$$

If  $x_A = A$  Alice performs U. If  $x_B = C$  Bob performs  $U^{-1}$ . Alice and Bob measure the state and output the resulting bit.

If  $x_A = B$  and  $x_B = B$  then no unitary transformation is performed and the players measure

$$\left| \Box \right\rangle = \frac{1}{\sqrt{2}} \left( 00 \right)_{AB} + \left| 11 \right\rangle_{AB} \right)$$

and obtain the same result with probability 1.

If  $x_A = A$  and  $x_B = B$  then Alice performs U on her qubit.

$$U \quad I|\Box\rangle = \frac{1}{\sqrt{2}} \left( U \quad I|00\rangle_{AB} + U \quad I|11\rangle_{AB} \right)$$

$$= \frac{1}{\sqrt{2}} \left( \frac{1}{\sqrt{2}} |0\rangle_{A} + \frac{1}{2} |1\rangle_{A} \right) \left( \frac{1}{2} |0\rangle_{A} + \frac{1}{2} |1\rangle_{A} \right) \left( \frac{1}{2} |0\rangle_{A} + \frac{1}{2} |1\rangle_{A} \right) \left( \frac{1}{2} |0\rangle_{A} + \frac{1}{2} |1\rangle_{A} \right)$$

$$= \frac{\sqrt{3}}{2\sqrt{2}} |00\rangle_{AB} + \frac{1}{2\sqrt{2}} |10\rangle_{AB} + \frac{1}{2\sqrt{2}} |01\rangle_{AB} + \frac{\sqrt{3}}{2\sqrt{2}} |11\rangle_{AB}$$

$$\Pr[y_A = y_B] = \begin{bmatrix} \sqrt{3} & 2 \\ 2\sqrt{2} & 2 \end{bmatrix} + \begin{bmatrix} \sqrt{3} & 2 \\ 2\sqrt{2} & 2 \end{bmatrix} = \frac{3}{4}$$

If  $x_A = B$  and  $x_B = C$  then Bob perform  $U^{-1}$  on his qubit.

$$I \quad U^{\square 1} | \square \rangle = \frac{1}{\sqrt{2}} \left( I \quad U^{\square 1} | 00 \rangle_{AB} + I \quad U^{\square 1} | 11 \rangle_{AB} \right)$$

$$= \frac{1}{\sqrt{2}} \left[ 0 \rangle_{A} \right] \frac{\sqrt{3}}{2} | 0 \rangle_{B} + \frac{1}{2} | 1 \rangle_{B} \left[ + | 1 \rangle_{B} \right] \frac{1}{2} | 0 \rangle_{A} + \frac{\sqrt{3}}{2} | 1 \rangle_{A} \left[ - \frac{\sqrt{3}}{2\sqrt{2}} | 00 \rangle_{AB} + \frac{1}{2\sqrt{2}} | 01 \rangle_{AB} + \frac{1}{2\sqrt{2}} | 10 \rangle_{AB} + \frac{\sqrt{3}}{2\sqrt{2}} | 11 \rangle_{AB} \right]$$

$$\Pr[y_A = y_B] = \begin{bmatrix} \sqrt{3} & 1 \\ 2\sqrt{2} & 1 \end{bmatrix} + \begin{bmatrix} \sqrt{3} & 1 \\ 2\sqrt{2} & 1 \end{bmatrix} = \frac{3}{4}$$

If  $x_A = A$  and  $x_B = C$  then Alice performs U on her qubit and Bob performs  $U^{-1}$  on his.

$$U \quad U^{\square} | \square \rangle = \frac{1}{\sqrt{2}} \left( U \quad U^{\square} | 00 \rangle_{AB} + U \quad U^{\square} | 11 \rangle_{AB} \right)$$

$$= \frac{1}{\sqrt{2}} \left( \frac{1}{\sqrt{2}} | 0 \rangle_{A} + \frac{1}{2} | 1 \rangle_{A} \right) \left( \frac{\sqrt{3}}{2} | 0 \rangle_{B} + \frac{1}{2} | 1 \rangle_{B} \right) + \left( \frac{1}{2} | 1 \rangle_{A} \right) \left( \frac{\sqrt{3}}{2} | 1 \rangle_{A} \right) \left( \frac{\sqrt{3}}{2} | 1 \rangle_{A} + \frac{\sqrt{3}}{2} | 1 \rangle_{A} \right)$$

$$= \frac{1}{2\sqrt{2}} | 00 \rangle_{AB} + \frac{\sqrt{3}}{2\sqrt{2}} | 01 \rangle_{AB} \left( \frac{\sqrt{3}}{4\sqrt{2}} | 10 \rangle_{AB} \right) \left( \frac{1}{2\sqrt{2}} | 11 \rangle_{AB} \right)$$

$$\Pr[y_A = y_B] = \frac{1}{2\sqrt{2}} = \frac{1}{2\sqrt{2}} = \frac{1}{2\sqrt{2}} = \frac{1}{4}$$

#### Interpretation

With very high probability Alice and Bob will violate the inequality that we have derived. This will enable them to win the game and achieve the pseudo-telepathy trick.

Alice and Bob are not strictly speaking exchanging information. Neither one learns the other player's input.

This proves that God plays dice and that quantum mechanics is not a local theory.

#### Mermin's game

	Alice	Bob	Charlie
Questions (1 bits)	$\mathcal{X}_{A}$	$\mathcal{X}_{B}$	$\mathcal{X}_{C}$
Answers (1 bits)	${\cal Y}_A$	${\cal Y}_B$	${\mathcal Y}_C$

Promise:  $x_A \oplus x_B \oplus x_C = 1$ 

Winning condition:  $y_A \oplus y_B \oplus y_C = x_A \square x_B \square x_C$ 

$$(x_A = 0, x_B = 0, x_C = 1)$$
  
 $(x_A = 0, x_B = 1, x_C = 0)$   
 $(x_A = 1, x_B = 0, x_C = 0)$   
Even  $(x_A = 1, x_B = 1, x_C = 1)$  Odd

#### Classical strategy?

Each player must have a strategy.

If 
$$x_A = 0$$
 Alice outputs  $y_A = a_0$   
If  $x_A = 1$  Alice outputs  $y_A = a_1$ 

If 
$$x_B = 0$$
 Bob outputs  $y_B = b_0$   
If  $x_B = 1$  Bob outputs  $y_B = b_1$ 

If 
$$x_C = 0$$
 Charlie outputs  $y_C = c_0$   
If  $x_C = 1$  Charlie outputs  $y_C = c_1$ 

Alice, Bob and Charlie's strategies must be consistent with the game's constraints.

#### No classical strategy!

$$x_A = 0, x_B = 0$$
 and  $x_C = 1$  
$$\begin{bmatrix} 0 = x_A & \exists x_B & \exists x_C \\ 0 = y_A \oplus y_B \oplus y_C \\ 0 = a_0 \oplus b_0 \oplus c_1 \end{bmatrix}$$
$$x_A = 0, x_B = 1 \text{ and } x_C = 0 \qquad \begin{bmatrix} 0 = a_0 \oplus b_1 \oplus c_0 \\ 0 = a_1 \oplus b_0 \oplus c_0 \end{bmatrix}$$
$$x_A = 1, x_B = 0 \text{ and } x_C = 0 \qquad \begin{bmatrix} 0 = a_1 \oplus b_0 \oplus c_0 \\ 1 = a_1 \oplus b_1 \oplus c_1 \end{bmatrix}$$

$$0 \oplus 0 \oplus 0 \oplus 1 = (a_0 \oplus b_0 \oplus c_1) \oplus (a_0 \oplus b_1 \oplus c_0) \oplus (a_1 \oplus b_0 \oplus c_0) \oplus (a_1 \oplus b_1 \oplus c_1)$$
$$1 = (a_0 \oplus a_0) \oplus (a_1 \oplus a_1) \oplus (b_0 \oplus b_0) \oplus (b_1 \oplus b_1) \oplus (c_0 \oplus c_0) \oplus (c_1 \oplus c_1)$$

$$1 = 0 \oplus 0 \oplus 0 \oplus 0 \oplus 0 \oplus 0$$

$$1 = 0$$

#### Classical players

Any classical strategy will be inconsistent with one of the four equations and therefor if the questions are chosen uniformly at random 25% of the time the players will fail the test.

Now suppose the players are separated by many lightyears but still succeed one million times in a row...

What should we conclude? Pseudo-telepathy???

#### Quantum players

$$H|0\rangle = \frac{1}{\sqrt{2}} (0) + |1\rangle$$

$$H|1\rangle = \frac{1}{\sqrt{2}} (0) |1\rangle$$

$$H(\Box|0\rangle + \Box|1\rangle) = \Box H|0\rangle + \Box H|1\rangle$$

$$= \Box \frac{|0\rangle + |1\rangle}{\sqrt{2}} + \Box \frac{|0\rangle \Box|1\rangle}{\sqrt{2}}$$

$$= \frac{\Box + \Box}{\sqrt{2}}|0\rangle + \frac{\Box\Box\Box}{\sqrt{2}}|1\rangle$$

$$HH\left(\square | 0\rangle + \square | 1\rangle\right) = H\left(\square + \square | 1\rangle\right) + \left(\square \square \square | 1\rangle\right) + \left(\square \square \square | 1\rangle\right) + \left(\square \square \square | 1\rangle\right)$$

$$HH = I$$
  $H = H^{\square 1}$ 

$$\left| \Box \right\rangle = \frac{1}{\sqrt{8}} \left( 000 \right\rangle + \left| 001 \right\rangle + \left| 010 \right\rangle + \left| 100 \right\rangle \Box \left| 011 \right\rangle \Box \left| 101 \right\rangle \Box \left| 110 \right\rangle \Box \left| 111 \right\rangle \right)$$

$$(H \quad I \quad I) |\Box\rangle = \frac{1}{2} (000) \Box |011\rangle + |101\rangle + |110\rangle) \Box$$

$$(I \quad H \quad I) |\Box\rangle = \frac{1}{2} (000) + |011\rangle \Box |101\rangle + |110\rangle) \Box$$

$$(I \quad I \quad H) |\Box\rangle = \frac{1}{2} (000) + |011\rangle + |101\rangle \Box |110\rangle) \Box$$

$$(H \quad H \quad H) |\Box\rangle = \frac{1}{2} (001) + |010\rangle + |100\rangle \Box |111\rangle) \quad \text{Odd}$$

The three players share the state  $\square$  Each player applies H if his input is 1. Each player measures his qubit and outputs the resulting bit.

$$\begin{aligned} & = \frac{1}{\sqrt{8}} \begin{bmatrix} (H & I & I) | 000 \rangle + (H & I & I) | 001 \rangle + (H & I & I) | 010 \rangle + (H & I & I) | 100 \rangle \\ & = \frac{1}{\sqrt{8}} \begin{bmatrix} (H & I & I) | 010 \rangle + (H & I & I) | 101 \rangle | | (H & I & I) | 110 \rangle | | (H & I & I) | 110 \rangle | | (H & I & I) | 111 \rangle \\ & = \frac{1}{\sqrt{8}} \begin{bmatrix} \frac{1}{\sqrt{2}} (0) + |1\rangle | 000 \rangle + \frac{1}{\sqrt{2}} (0) + |1\rangle | 01 \rangle + \frac{1}{\sqrt{2}} (0) + |1\rangle | 10 \rangle + \frac{1}{\sqrt{2}} (0) + |1\rangle | 10 \rangle | | \frac{1}{\sqrt{2}} (0) + |1\rangle | 110 \rangle | | |111 \rangle | |$$

#### Asymptotic result

There exists a pseudo-telepathy game that the players can win with probability 1 if they share *n* Bell states.

Without entanglement this game requires  $\prod (2^n)$  bits of communication.

This gives an exponential bound on the amount of communication required to simulate *n* entangled qubits.